



TITLE: RESOURCE MANAGEMENT APPARATUS, SYSTEMS, AND METHODS

INVENTOR'S NAME: SACHIN DOSHI, ET AL.

SERIAL NO.: 10/750,933 DOCKET NO.: 884.A60US1

REPLACEMENT SHEET

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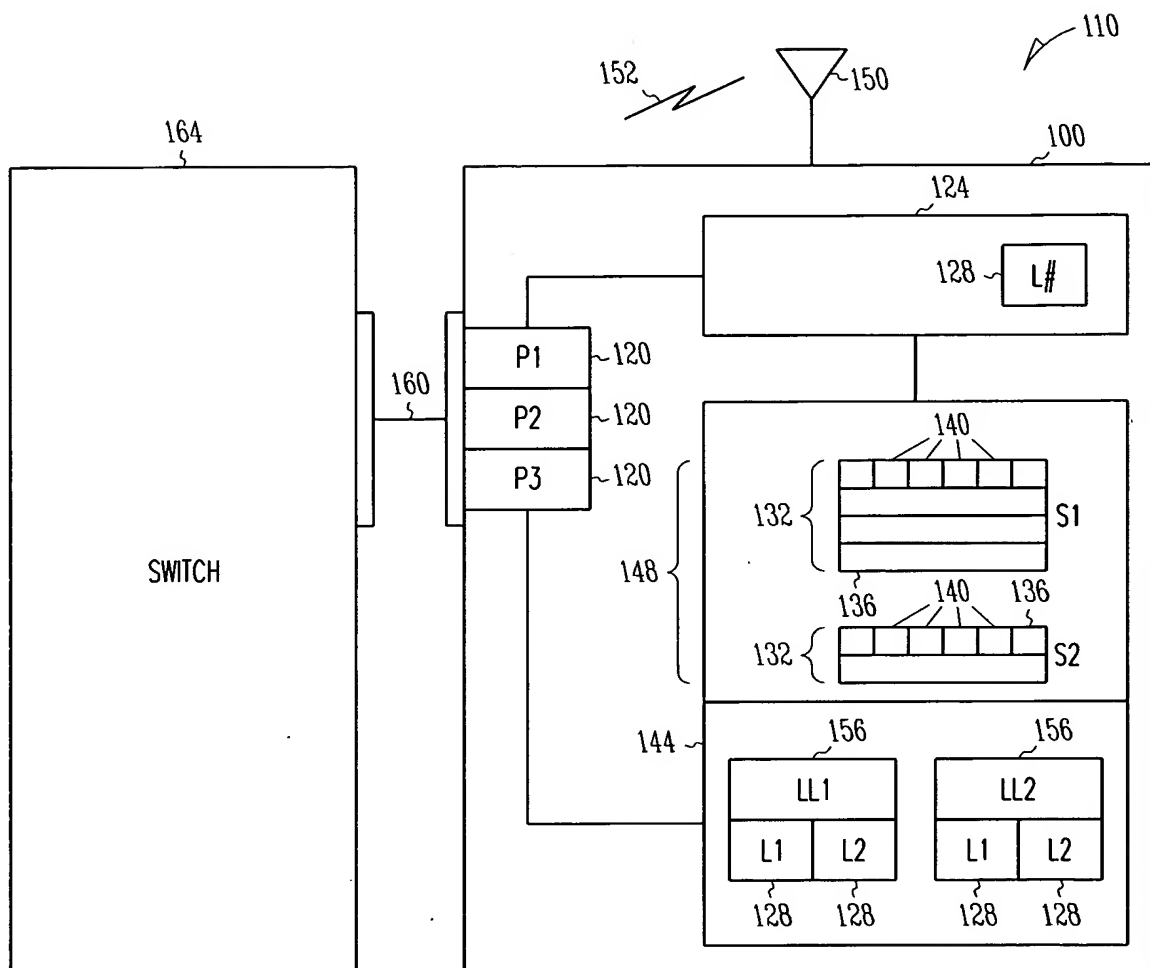


FIG. 1

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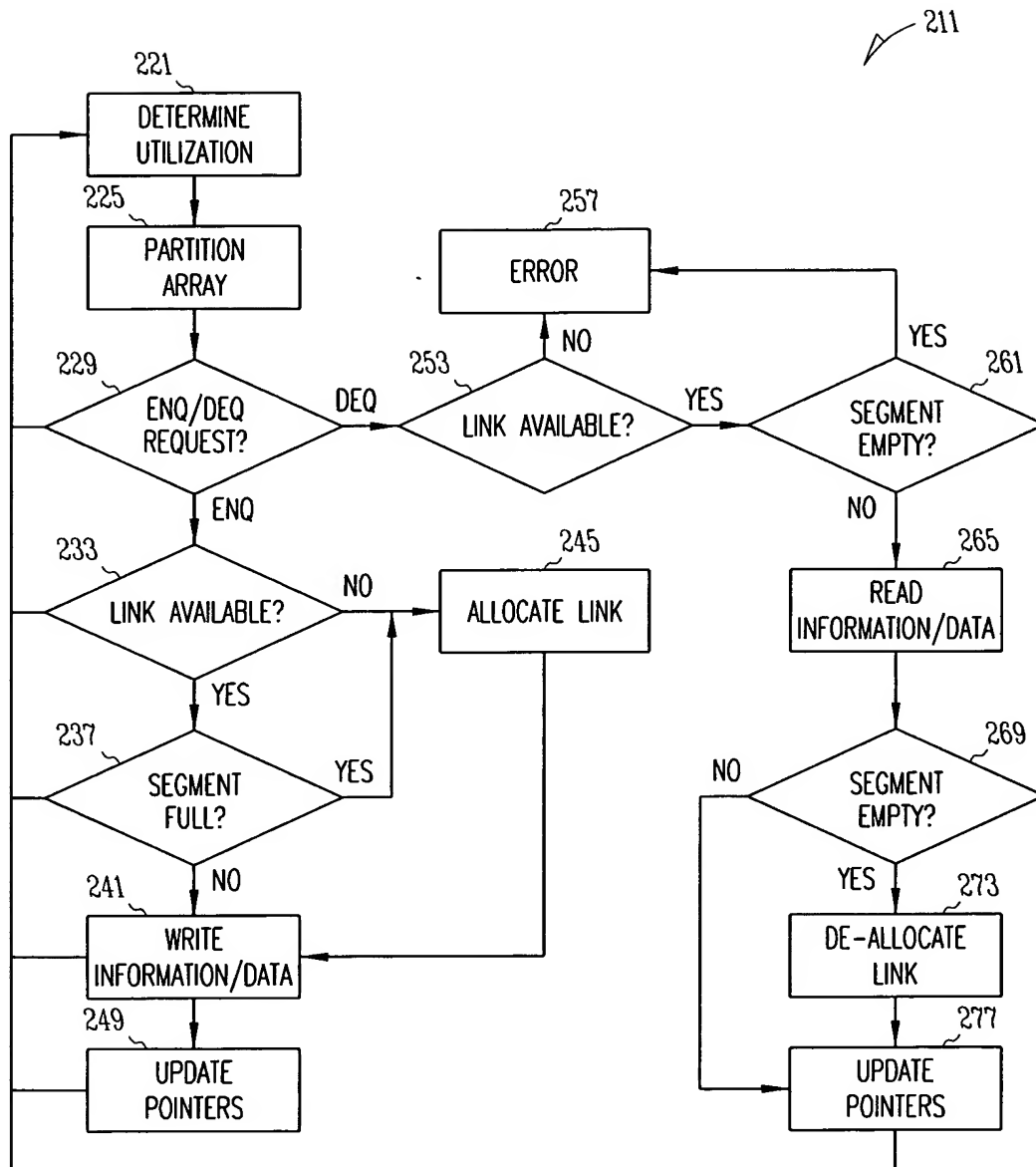


FIG. 2

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ENQ_OPERATION

```

// This is pseudo code for adding an entry into the link list.
// WP - Write pointer
// Info_array: The information array
// cur_port_id: the ID of the port for which the enqueue operation
// is being done. Note that many ports are sharing this
// resource.
// LR - Link RAM, linked list information stored here
// This pseudo code is for a 4 to 1 linked list
}
If (WP[cur_port_id][1:0]==2'b00){~373
// If the lower 2 bits of the WP of the current port is 0, a new segment
// is required for this port.
// Take the free_avail_link that provides a pointer to the first free segment
// and make the WP point to the next location (01) in the new segment.
// Store the information in the 1st location (00) of the new segment.
WP[cur_port_id]<={free_avail_link, 2'b01};
Info array [{free_avail_link, 2'b00}] <= Info; ~375
} Else {
// If the lower 2 bits of the WP are non-zero, it means that the segment has
// space to store some more information. Keep adding the information in the
// empty locations of the segment. Note that the information storage is
// sequential within a segment.
WP[cur_port_id] <= {WP[cur_port_id][11:2], WP[cur_port_id][1:0]+2'b01}; ~377
Info_array[WP[cur_port_id]]<= Info;
}

// This portion keeps track of whether the link list for a port is empty
// or not. It also updates the linked list with the new segments.
if ( Empty_Flag[cur_port_id])
{ if (WP[cur_port_id][1:0]==2'b00)
LR[WP[cur_port_id][11:2]]<=free_avail_link; ~379
// If the linked list for the current port is not empty then whenever
// a new segment is allocated, link it in the link ram of the current
// port.
}
Else
// If the linked list is empty for a particular port then initialize the
// read and write pointers. Also reset the empty flag.
// Note that the read pointer will be updated by the dequeue operation.
{ Empty_Flag[cur_port_id] <= FALSE; ~381
if (WP[cur_port_id][1:0]==2'b00)
RP[cur_port_id] <={avail_link_S0, 2'b00};
}
}

```

DEQ_OPERATION

```

// this is the pseudo code for the dequeue operation.

Cur_rp = RP[cur_port_id]; ~383
// generation of the empty condition for the dequeue operation
Empty_condition = WP[cur_port_id]=={Cur_rp[11:2], Cur_rp[1:0]+2'b01}
If (~Empty_condition && Cur_rp[1:0]==2'b11) ~385
// If the linked list is not empty and we are reading the last location
// within a segment, then take the read pointer from the link ram.
RP[cur_port_id] <= {LR[RP[cur_port_id][11:2]], 2'b00};
} else {
// If we are not reading the last location within a segment then keep on
// incrementing the read pointer within the segment.
RP[cur_port_id] <= {RP[cur_port_id][11:2], RP[cur_port_id][1:0] + 2'b01}; ~387
}

// When the segment is completely read, put the free segment in the pool of the
// free segments.
If(RP[cur_port_id][1:0]==2'b11) ~389
Put_free_link (RP[qnum_s0][11:2]); //LR write

// Set the empty flag whenever the empty condition is detected for a particular
// port
If (Empty_condition) Empty_Flag[cur_port_id] = TRUE; ~391

```

FIG. 3

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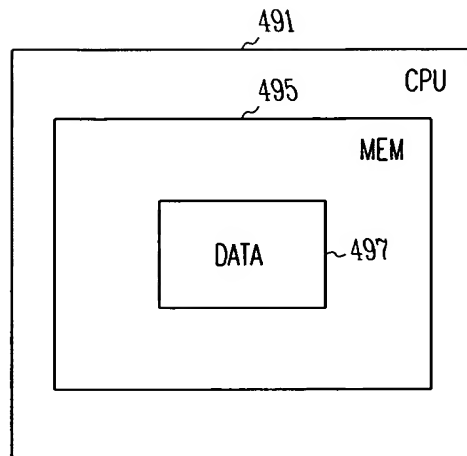


FIG. 4